



Data structure and Algorithm

By
Elsayed Atlam
2022

Course Contents

- Data Types
 - Overview, Introductory concepts
 - Data Types, meaning and implementation
 - Abstract data types (ADT)
 - Arrays (revisited)
 - Structures
- Stacks (recursion)
- Queues
- Linked Lists
- Trees (traversals, implementation)

Course Contents

- Binary Trees
- Indexing Methods
 - Hashing
- Binary Search Trees
- Balanced Search Trees
 - (AVL Tree) Adelson-Velskii-Landis
- Heaps

Course objectives

- ❑ Be familiar with different data structures available to represents data
- ❑ Be able to trace algorithms and verify correctness.
- ❑ Be able to develop and implement algorithms using different data structures
- ❑ Be able to select appropriate data structures and algorithms for given problems
- ❑ Be able to use JAVA language to implement different algorithms pseudo codes.

Objectives of the course

- Present in a systematic fashion the most commonly used data structures, emphasizing their abstract properties.
- Discuss typical algorithms that operate each kind of data structure, and analyze their performance.
- Compare different Data Structures for solving the same problem, and choose the best

Readings/references

□ Text Book:

- **Data Structures & Algorithms in JAVA** (5th Edition), by M. Goodrich & R. Tamassia, John Wiley & Sons, inc., 2010.

□ Additional Readings:

- **Data Structures and Problem Solving with JAVA** (3rd Edition), by Mark Allen Weiss, Addison Wesley, 2006.
- **Lecture slides and handouts**

What is data?

□ Data

- A collection of facts from which conclusion may be drawn
- e.g. **Data**: Temperature 35°C; **Conclusion**: It is hot.

□ Types of data

- Textual: For example, your name (Muhammad)
- Numeric: For example, your ID (090254)
- Audio: For example, your voice
- Video: For example, your voice and picture
- (...)

What is the difference between Data and Information?

- Data are a set of collected numbers, words, anything. They do not mean anything until they are organized, arranged or developed.
- **Examples:** *numbers, dates, prices, names (Olive)*
 - Once that happens (after they have been **processed**), information is obtained.
- Information actually makes sense and is expressed through some sort of comprehensible logic.
- **Examples:** *reports, Tables, Figures (olive Oil)*

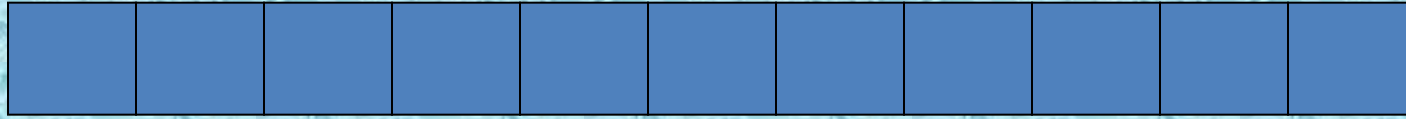
What is the Processing that change Data to Information?

- Adding
- Deleting
- Multiplying
- Logical operations: !=, ==, etc
- Retrieving, modifying, updating, saving

What is data structure?

- ❑ A particular way of **storing and organizing data in a computer** so that it can be used efficiently and effectively.
- ❑ Data Structures are **the programmatic way of storing data so that data can be used efficiently**.
- ❑ Data structure is the logical or mathematical model of a particular organization of data.
- ❑ A group of data elements grouped together under one name.
 - For example, an array of integers

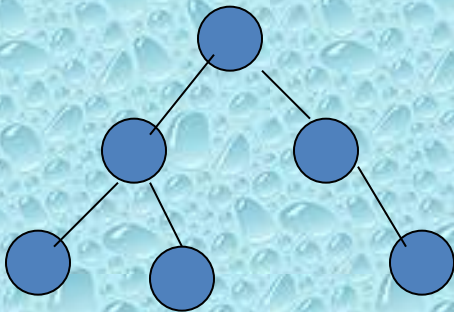
Types of data structures



1-D Array



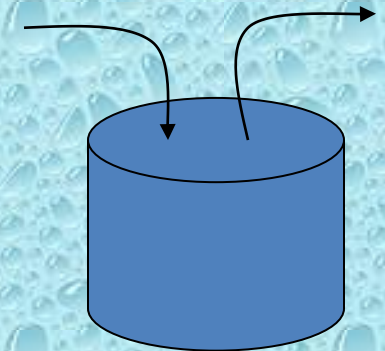
Linked List



Tree



Queue



Stack

There are many, but we named a few. We'll learn these data structures in great detail!

The Need for Data Structures

- ❑ **Goal:** to **organize data**

- ❑ **Criteria:** to facilitate **efficient**
 - **storage** of data
 - **retrieval** of data
 - **manipulation** of data

- ❑ **Design Issue:**
 - **select and design** appropriate data types
(This is the main motivation to learn and understand data structures)

Why Study?

- A particular way of **storing and organizing data in a computer** so that it can be used efficiently and effectively
- Designed to develop students understanding the impact of *structuring data to achieve efficiency* of a solution to a problem
- After completion you will be familiar with important and most often used data structuring techniques.
- It will enable you to understand the manner in which data is organized and presented later.

Data Structure Operations

(Demonstrate using class room example!)

❑ Traversing

- Accessing each data element exactly once so that certain items in the data may be processed

❑ Searching

- Finding the location of the data element (key) in the structure

❑ Insertion

- Adding a new data element to the structure

Data Structure Operations (cont.)

❑ Deletion

- Removing a data element from the structure

❑ Sorting

- Arrange the data elements in a logical order (ascending/descending)

❑ Merging

- Combining data elements from two or more data structures into one

What is algorithm?

- ❑ A finite set of instructions which accomplish a particular task
- ❑ A method or process to solve a problem
- ❑ Transforms input of a problem to output

Algorithm = Input + Process + Output

Algorithm development is an art – it needs practice, practice and only practice!

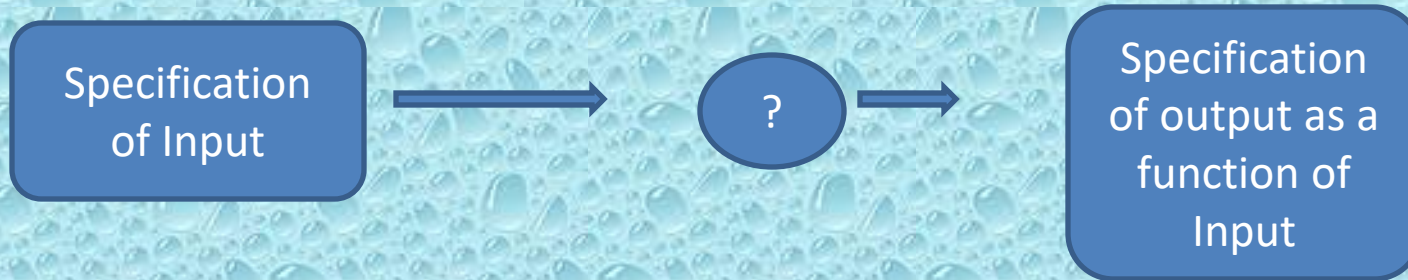
- Algorithm is a step-by-step procedure, which defines a set of instructions to be executed in a certain order to get the desired output.
- Algorithms are generally created independent of underlying languages, i.e. an algorithm can be implemented in more than one programming language.
- From the data structure point of view, following are some important categories of algorithms –
 - **Search** – Algorithm to search an item in a data structure.
 - **Sort** – Algorithm to sort items in a certain order.
 - **Insert** – Algorithm to insert item in a data structure.
 - **Update** – Algorithm to update an existing item in a data structure.
 - **Delete** – Algorithm to delete an existing item from a data structure.

Introduction

Data structure and Algorithm

- ❑ Algorithm: outline, the essence of a computational procedure, step by step instructions
- ❑ Program: an implementation of an Algorithm, written in some specific programming language
- ❑ Data Structure: Organization of Data needed to solve the problem

Algorithmic Problem

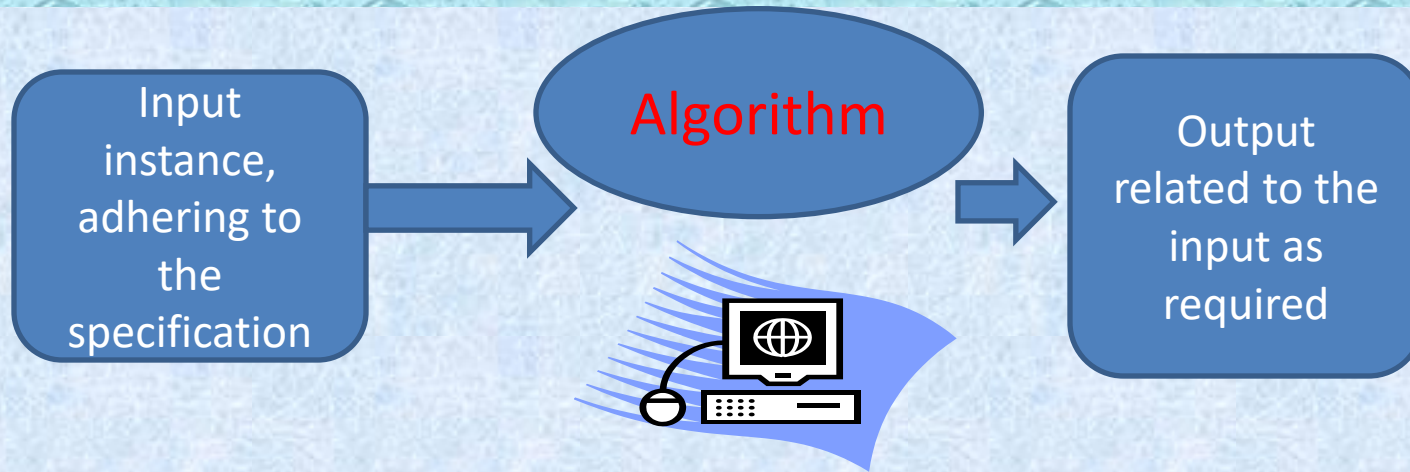


□ Infinite number of input instances satisfying the specification, For example : A sorted, non-decreasing sequence of natural number of non-zero, finite length:

□ 1, 20, 908, 909, 100000, 1000000

□ 3

Algorithmic Solution



- Algorithm describes actions on the input instance to get an output as desired as specified
- Again infinitely many correct algorithms can be used for the same Algorithmic problem

What is a good Algorithm?

- ❑ Efficient: Any thing is efficient is good
 - ❑ Small Running time
 - ❑ Space Used (Less Memory)

What is a good algorithm?

- ❑ It must be correct
- ❑ It must be finite (in terms of time and size)
- ❑ It must terminate
- ❑ It must be unambiguous
 - Which step is next?
- ❑ It must be space and time efficient

A program is an instance of an algorithm, written in some specific programming language

What is a good Program?

There are a number of facets to good programs: they must

- run correctly
- run efficiently
- be easy to read and understand
- be easy to debug *and*
- be easy to modify.

What does correct mean?

We need to have some formal notion of the meaning of correct: thus we define it to mean

"run in accordance with the specifications".

A simple algorithm

❑ **Problem:** Find maximum of a , b , c

❑ **Algorithm**

▪ Input = a , b , c

▪ Output = max

▪ Process

○ Let $max = a$

○ If $b > max$ then
 $max = b$

○ If $c > max$ then
 $max = c$

○ Display max

Order is very important!!!

Algorithm development: Basics

□ Clearly identify:

- what output is required?
- what is the input?
- What steps are required to transform input into output
 - The most crucial bit
 - Needs problem solving skills
 - A problem can be solved in many different ways
 - Which solution, amongst the different possible solutions is optimal?

How to express an algorithm?

- ❑ A sequence of steps to solve a problem
- ❑ We need a way to express this sequence of steps
 - **Natural language** (NL) is an obvious choice, but not a good choice. Why?
 - NLs are notoriously ambiguous (unclear)
 - **Programming language** (PL) is another choice, but again not a good choice. Why?
 - Algorithm should be PL independent
 - **We need some balance**
 - We need PL independence
 - We need clarity
 - Pseudo-code provides the right balance

What is pseudo-code?

- ❑ Pseudo-code is a short hand way of describing a computer program
- ❑ Rather than using the specific syntax of a computer language, more general wording is used
- ❑ It is a mixture of NL and PL expressions, in a systematic way
- ❑ Using pseudo-code, it is easier for a non-programmer to understand the general workings of the program

Pseudo-code: general guidelines

- ❑ **Use PLs construct** that are consistent with modern high level languages, e.g. C++, Java, ...
- ❑ **Use appropriate comments** for clarity
- ❑ **Be simple and precise**

Pseudo-Code

□ A mixture of natural language and high –level programming concepts that describes the main idea behind a generic implementation of a data structure or Algorithm.

□ Eg: **Algorithm arrayMax(A,n):**

Input: An array A storing n integers,

Output: the maximum element in A.

currentMax \leftarrow A[0]

for i \leftarrow 1 to n-1 do

if currentMax $<$ A[i] then currentMax \leftarrow A[i]

return currentMax

Pseudo-Code

It is **more structured** than usual prose but **less formal** than a programming language

What pseudo-code looks like:

Expressions:

- Use standard mathematical symbols to describe numeric and boolean expressions
- Use \leftarrow for assignment (`=` in C)
- Use `=` for the equality relationship (`==` in C)

Method Declaration

- Algorithm name (param1, Param2)

Pseudo-Code

- ❑ Programming Constructions:
 - ❑ Decision structure: **if... then... [else....]**
 - ❑ While-loops: **while....do**
 - ❑ Repeat-loops: **repeat.... Until....**
 - ❑ For-loop: **for....do**
 - ❑ Array indexing: **A[i], A[l,j]**

- ❑ Methods
 - ❑ Calls: object method(args)
 - ❑ Returns: **return** Value

Components of Pseudo-code With Examples

□ Expressions

- Standard mathematical symbols are used
 - Left arrow sign (\leftarrow) as the **assignment operator** in assignment statements (equivalent to the = operator in Java)
 - Equal sign ($=$) as the **equality relation** in Boolean expressions (equivalent to the "= =" relation in Java)
 - For example

Sum \leftarrow 0

Sum \leftarrow Sum + 5

What is the final value of sum?

Components of Pseudo-code (cont.)

- ❑ **Decision structures** (if-then-else logic)
 - **if** condition **then** true-actions [**else** false-actions]
 - We use indentation to indicate what actions should be included in the true-actions and false-actions
 - For example

```
if marks > 50 then  
    print "Congratulation, you are passed!"  
else  
    print "Sorry, you are failed!"  
end if
```

What will be the output if marks are equal to 75?

Components of Pseudo-code (cont.)

□ Loops (Repetition)

▪ Pre-condition loops

○ While loops

- **while** condition **do** actions

- We use indentation to indicate what actions should be included in the loop actions

- For example

```
while counter < 5 do  
    print “Welcome to CS204!”  
    counter ← counter + 1  
end while
```

What will be the output if counter is initialised to 0, 7?

Components of Pseudo-code (cont.)

□ Loops (Repetition)

▪ **Pre-condition loops**

○ For loops

- **for** variable-increment-definition **do**
actions

- For example

```
for counter ← 0; counter < 5; counter ←  
counter + 2 do  
    print “Welcome to CS204!”  
end for
```

What will be the output?

Components of Pseudo-code (cont.)

□ Loops (Repetition)

▪ **Post-condition loops**

○ Do loops

- **do** actions **while** condition
- For example

do

print “Welcome to CS204!”

counter ← counter + 1

while counter < 5

What will be the output, if counter was initialised to 10?

The body of a post-condition loop must execute at least once

Homework

1. Write an algorithm to find the largest of a set of 10 numbers.
2. Write an algorithm in pseudocode that finds the average of (10) numbers.

Components of Pseudo-code (cont.)

□ Method declarations

- **Return_type** **method_name** (**parameter_list**)
 method_body

- For example

```
integer sum ( integer num1, integer num2)
```

```
    start
```

```
        result ← num1 + num2
```

```
    end
```

□ Method calls

- **object.method (args)**
- For example

```
    mycalculator.sum(num1, num2)
```

Components of Pseudo-code (cont.)

□ Method returns

- **return** value

- For example

```
integer sum ( integer num1, integer  
num2)
```

```
start
```

```
    result ← num1 + num2
```

```
    return result
```

```
end
```

Components of Pseudo-code (cont.)

□ Comments

- `/* Multiple line comments go here. */`
- `// Single line comments go here`
- Some people prefer braces `{}`, for comments

□ Arrays

- $A[i]$ represents the i th cell in the array A .
- The cells of an n -celled array A are indexed from $A[0]$ to $A[n - 1]$ (consistent with Java).

Algorithm Design: Practice

- Example : Determining even/odd number
 - A number divisible by 2 is considered an even number, while a number which is not divisible by 2 is considered an odd number. Write pseudo-code to display first N odd/even numbers.

Even/ Odd Numbers

Input range

```
for      num←0;          num<=range;  
  num←num+1 do  
  if num % 2 = 0 then  
    print num is even  
  else  
    print num is odd  
  endif  
endfor
```

1. Write an algorithm to find the largest of a set of 10 numbers.

Input: 10 positive integers

Output: Max integer

Process:

Range=10;

Max ← 0;

Counter ← 1;

```
for counter ← 0; counter ≤ range;
  counter ← counter + 1 do
  if integer ≥ max then
    max = integer;
  endif
```

Endfor

Return max;

FindLargest

Input: 1000 positive integers

1. Set Largest to 0
 2. Set Counter to 0
 3. while (Counter less than 1000)
 - 3.1 if (the integer is greater than Largest)
then
 - 3.1.1 Set Largest to the value of the integer
 - End if
 - 3.2 Increment Counter
 - End while
 4. Return Largest
- End**

1. Write an algorithm in pseudocode that finds the average of (10) numbers.

Input: 10 positive integers

Output: average of 10 integers

Process:

sum \leftarrow 0;

for $i \leftarrow 0$; $i \leq 10$; $i \leftarrow i + 1$ do

 input x;

 sum = sum + x;

Endfor

Avg = sum / 10;

Return Avg;

Write an algorithm which requires a number between 10 and 20, until the response is appropriate. If the number is more than 20, it will display a message: "Bigger!" If the number is less than 10, it will display "smaller!"

Begin

Input: num

Output: numbers between 10 and 20

Process:

Start

if (num < 10) Then

print "Smaller !"

elseif (num > 20)

print "Bigger !"

End if

End

What are the values of the variables A, B and C after execution of the following instructions?

Begin

$A \leftarrow 3$

$B \leftarrow 10$

$C \leftarrow A + B$

$B \leftarrow A + B$

$A \leftarrow C$

End

Write an algorithm to swap the value the 2 variables A and B.

Input: A and B and C

Output: Swapping

Process:

Start

C ← A;

A ← B;

B ← C;

Return A and B;

End

Write pseudocode that will take a number as input and tells whether a number is positive, negative or zero.

Begin

WRITE "Enter a number"

READ num

IF num > 0 THEN

WRITE "The number is positive"

ELSE IF num = 0 THEN

WRITE "The number is zero"

ELSE

WRITE "The number is negative"

ENDIF

ENDIF

End

Question?

- Write a pseudo-code to count (calculate) the submission of the first 100 normal number?
- what is a good Algorithm?
- What is a good program?

Measuring the Running time

How should we measure the running time of an **Algorithm**?

Experimental Study

- ❑ Write a certain **program** that implements the algorithm
- ❑ Run the program with data sets of varying size (large or small) and composition
- ❑ Clock the time by: Use a method like **System.currentTimeMillis()** to get an accurate measure of the actual running time

Limitations of Experimental Studies

- ❑ It is necessary to **implement** and test the algorithm in order to determine its running time.
- ❑ Experiments can be done only on **a limited set of inputs**, and may not be indicative of the running time on other inputs not included in the experiment.
- ❑ In order to compare two algorithms, the same **Hardware and software environments** should be used.

Beyond Experimental Studies

We will develop a **general methodology** for analyzing running time of algorithms. This approach (we want to)

- Uses a **high-level description** of the algorithm instead of testing one of its implementations.
- Takes into account **all possible inputs**
- Allows one to evaluate the efficiency of any algorithm in a way that is **independent of the hardware and software environment**

Analysis of Algorithms

❑ **Primitive Operation:** low-level operation independent of programming language.

Can be identified in pseudo-code. For eg:

- ❑ Data movement (assign)
- ❑ Control (Branch, subroutine call, return)
- ❑ Arithmetic and logical operations (e.g. addition, comparison)
- ❑ By inspecting the pseudo-code, we can count the number of primitive operations executed by the algorithm

Example: Sorting

INPUT

Sequence of numbers

$a_1, a_2, a_3, \dots, a_n$



2 5 4 10 7

Sort



$b_1, b_2, b_3, \dots, b_n$

2 4 5 7 10

OUTPUT

A permutation of the sequence of numbers

Correctness (requirements for the output)

For any given input the algorithm halts with the output

- $b_1 < b_2 < b_3 < \dots < b_n$
- b_1, b_2, \dots, b_n is a permutation of a_1, a_2, \dots, a_n

Running time Depend on

- Number of element (n)
- How (partially) sorted they are
- Algorithm

Insertion Sort



Cards Hand Play




```
void insertionSort(int arr[], int n)
{
    int key, j;
    for (int i = 1; i < n; i++)
    {
        key = arr[i];
        j = i - 1;
        while (j >= 0 && arr[j] > key)
        {
            arr[j + 1] = arr[j];
            j = j - 1;
        }
        arr[j + 1] = key;
    }
}
```

//0 1 2 3 4 5 6
//80 |90| 60 30 50 70 40
I

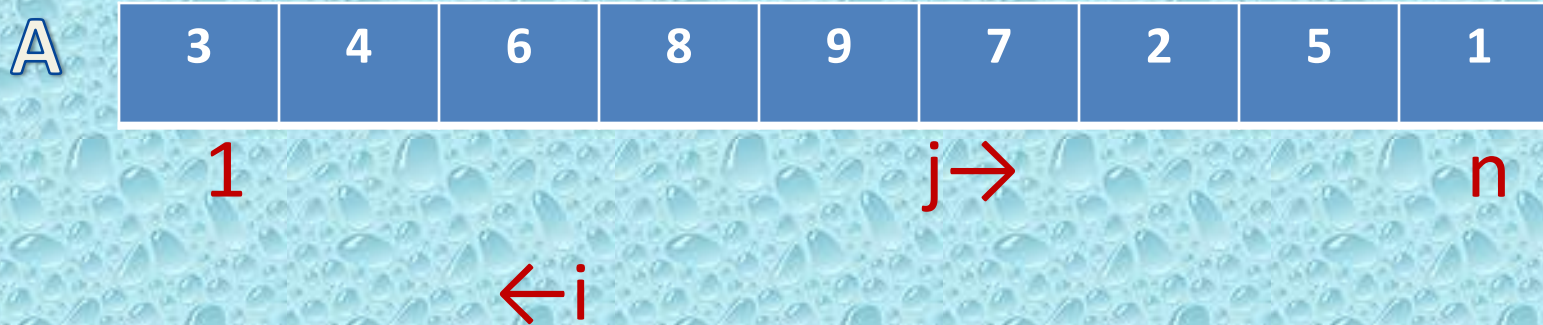
```

void insertionSort(int arr[], int n)
{
    int key, j; //0 1 2 3 4 5 6
    for (int i = 1; i < n; i++) //60 80 90 || 30 50 70 40
    {
        key = arr[i]; //60 i=2
        j = i - 1; //0

        while (j >= 0 && arr[j] > key)
        {
            arr[j + 1] = arr[j];
            j = j - 1;
        }
        arr[j + 1] = key; //90
    }
}

```

Example: Sorting



Strategy

- Start empty handed
- Insert a card in the right position of the already sorted hand
- Continue until all cards are inserted sorted

INPUT: $A[1\dots n]$ - an array of integers
OUTPUT: a permutation of A such that $A[1] < A[2] < \dots < A[n]$

```
for j=2 to n do
  Key  $\leftarrow$  A[j]
  Insert A[j] into the sorted sequence
  A[1,j-1]
  i  $\leftarrow$  j-1
  While i > 0 and A[i] > Key
    do A[i+1]  $\leftarrow$  A[i]
    i - -
  A[i+1]  $\leftarrow$  key
```

Analysis of Algorithms

Algorithm	Cost	Times
for j=2 to n do	C1	n-1
Key ← A[j]	C2	n-1
Insert A[j] into the sorted sequence A[1,j-1]	0	n-1
i ← j-1	C3	n-1
While i > 0 and A[i] > Key	C4	$\sum_{j=2}^n t_j$
do A[i+1] ← A[i]	C5	$\sum_{j=2}^n t_j - 1$
i - -	C6	$\sum_{j=2}^n t_j - 1$
A[i+1] ← key	C7	n-1

$$\text{Total time} = n(C1+C2+C3+C7) + \sum_{j=2}^n t_j (C4+C5+C6) - (C1+C2+C3+C5+C6+C7)$$

Best/Worst/average Case (1)

$$\text{Total time} = n(C1+C2+C3+C7) + \sum_{j=2}^n t_j (C4+C5+C6) - (C1+C2+C3+C5+C6+C7)$$

❑ **Best Case:** elements already sorted; $t_j=1$,

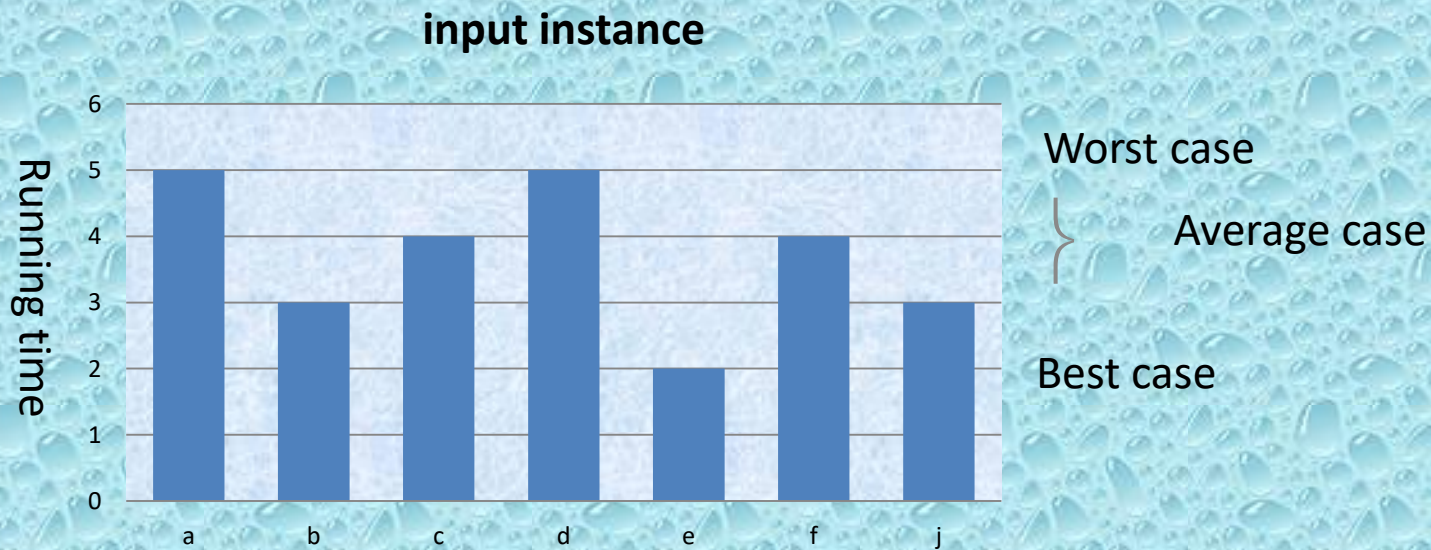
Running time = $f(n)$, i.e. Linear time

❑ **Worst Case:** elements are sorted in inverse order, $t_j=j$,
running time = $f(n^2)$, i.e quadratic time

❑ **Average case:** $t_j=j/2$, running time = $f(n^2)$
i.e quadratic time

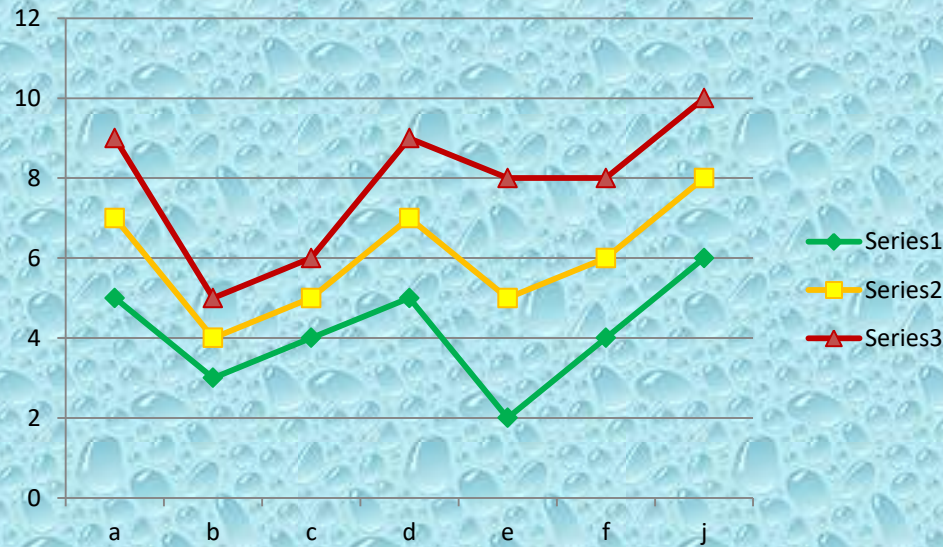
Best/Worst/average Case (2)

For a specific size of input n , investigate running times for different input instance



Best/Worst/average Case (3)

For inputs of all sizes:



Best/Worst/average Case (4)

- ❑ **Worst Case:** is usually used: it is an upper bound and in certain application domains (e.g. air traffic control, surgery) knowing the **worst case** time complexity is of crucial important.
- ❑ For some algorithms **worst case** occurs fairly often
- ❑ **Average case:** is often as bad as the **worst case**
- ❑ Finding **average case** can be very difficult

Asymptotic Analysis

- ❑ Goal: to simplify analysis of running time by getting rid of details which may be affected by specific implementation and hardware
 - ❑ Like *rounding*: $1,000,001 = 1,000,000$
 - ❑ $3n^2 = n^2$
- ❑ **Capturing the essence:** how the running time of an algorithm increases with the size of the input in the limit.
 - ❑ **Asymptotic** more efficient algorithms are best for all but small inputs

Asymptotic Analysis of Running time

- ❑ Using O -notation to express number of primitive operations executed as function of input size.
- ❑ Comparing asymptotic running times:
 - ❑ An Algorithm that runs in $O(n)$ is better than one runs in $O(n^2)$ time
 - ❑ Similarly , $O(\log n)$ is better than $O(n)$
 - ❑ Hierarchy of functions: $\log n < n < n^2 < n^3 < 2^n$
- ❑ **Caution!** Beware of very large constant factors. An algorithm running in time $1,000,000 n$ is still $O(n)$ but might be less efficient than one running in time $2n^2$ which is $O(n^2)$

Example of Asymptotic Analysis

Algorithm of prefix Averages1(X):

Input: An n-element array X of Numbers

Output: An n-element array A of numbers such that A[i] is the average of elements X[0],...,X[i]

for $i \leftarrow 0$ to $n-1$ do

$a \leftarrow 0$

for $j \leftarrow 0$ to i do

$a \leftarrow a + X[j]$

$A[j] \leftarrow a / (i+1)$

return array A

i iterations with
 $i=0,1, n-1$
Executed I times

n iterations

Analysis: running time is $O(n^2)$

Example of Asymptotic Analysis (A Better Algorithm)

Algorithm of prefix Averages₂(X):

Input: An n -element array X of Numbers

Output: An n -element array A of numbers such that $A[i]$ is the average of elements $X[0], \dots, X[i]$

$S \leftarrow 0$

for $i \leftarrow 0$ to n do

$S \leftarrow S + X[i]$

$A[i] \leftarrow S / (i + 1)$

return array A

Analysis: running time is $O(n)$

Example of Asymptotic Analysis (A Better Algorithm)

Algorithm of prefix Averages1(X):

Input: An n-element array X of Numbers

Output: An n-element array A of numbers such that A[i] is the average of elements X[0],...,X[i]

for $i \leftarrow 0$ to $n-1$ do

$a \leftarrow 0$

 for $j \leftarrow 0$ to i do

$a \leftarrow a + X[j]$

$A[j] \leftarrow a / (i+1)$

 return array A

Analysis: running time is $O(n^2)$

Algorithm of prefix Averages2(X):

Input: An n-element array X of Numbers

Output: An n-element array A of numbers such that A[i] is the average of elements X[0],...,X[i]

$S \leftarrow 0$

for $i \leftarrow 0$ to n do

$S \leftarrow S + X[i]$

$A[i] \leftarrow S / (i+1)$

 return array A

Analysis: running time is $O(n)$

Comparison of running Times

Running Time	Maximum Problem Size (n)		
	1 Second	1 minute	1 hour
$400n$	2500	150000	9000000
$20 n \log n$	4096	166666	7826087
$2n^2$	707	5477	42426
n^4	31	88	244
2^n	19	25	31

you can see what is the largest size of the problem you can solve in one second , 1 minutes and 1 hour,

Notice

- نلاحظ ان سرعة معالجة البيانات تعتمد علي عدة عوامل إضافة إلي العامل الزمني اللازم للمعالجة مثل:
- عوامل تحديد الذاكرة الرئيسية
- عوامل تحديد وحدات الإدخال والإخراج
- عوامل تحديد تفاعل وحدات الإدخال والإخراج مع الذاكرة الرئيسية
- What is Big-O notation?
- Explanation about Array?(delete, insert,...)